

## Matrix Multiplication

- Sequential algorithm
- $t(n) = O(n^3)$

**procedure MATRIX MULTIPLICATION ( $A, B, C$ )**

**for  $i = 1$  to  $m$  do**

**for  $j = 1$  to  $k$  do**

(1)  $c_{ij} \leftarrow 0$

(2) **for  $s = 1$  to  $n$  do**

$c_{ij} \leftarrow c_{ij} + (a_{is} \times b_{sj})$

**end for**

**end for**

**end for.  $\square$**

## Parallel Matrix Multiplication

On mesh connected computer

**procedure** MESH MATRIX MULTIPLICATION ( $A, B, C$ )

**for**  $i = 1$  to  $m$  **do in parallel**

**for**  $j = 1$  to  $k$  **do in parallel**

(1)  $c_{ij} \leftarrow 0$

(2) **while**  $P(i, j)$  receives two inputs  $a$  and  $b$  **do**

(i)  $c_{ij} \leftarrow c_{ij} + (a \times b)$

(ii) **if**  $i < m$  **then** send  $b$  to  $P(i + 1, j)$

**end if**

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(iii) **if**  $j < k$  **then** send  $a$  to  $P(i, j + 1)$

**end if**

**end while**

**end for**

**end for.**  $\square$

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Ex-

$$A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} \quad \text{and} \quad B = \begin{bmatrix} -5 & -6 \\ -7 & -8 \end{bmatrix}$$

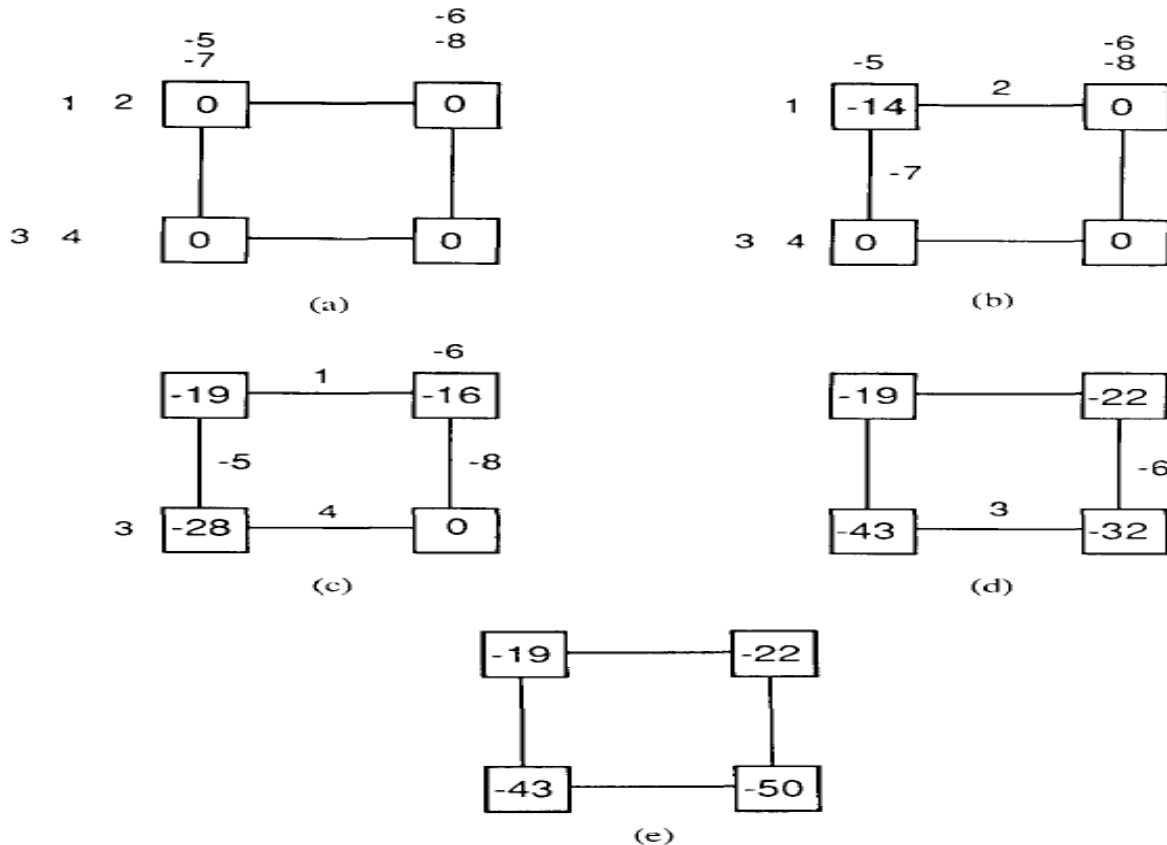


Figure 7.8 Multiplying two matrices using procedure MESH MATRIX

### Analysis-

- Elements of first row of A and first column of B takes  $(m+k+n-2)$  steps to reach the last processor, i.e.,  $P(m,k)$
- $t(n) = O(n)$ , as  $m \leq n$  and  $k \leq n$  (assumed)

- $c(n) = p(n) * t(n)$

$= n^2 * n = O(n^3)$ , i.e., cost-optimal

Another cost-optimal algorithm is CRCW MATRIX  
MULTIPLICATION