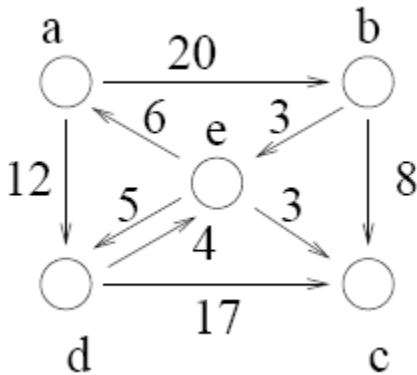
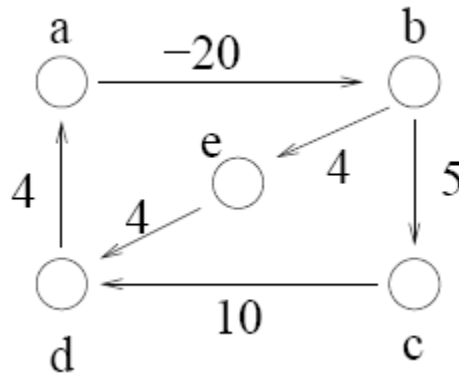


## The All-Pairs Shortest Paths Problem

Given a weighted digraph  $G(V,E)$  with weight function  $w: E \rightarrow \mathbb{R}$ , ( $\mathbb{R}$  is the set of real numbers), determine the length of the shortest path (i.e., distance) between all pairs of vertices in  $G$ . Here we assume that there are no cycles with zero or negative cost.



Without negative cost cycle



With negative cost cycle

**Solution1:** If there are no negative cost edges apply Dijkstra's algorithm to each vertex (as the source) of the digraph. Recall the Dijkstra's algorithm run in  $\mathcal{O}(V+E(\log V))$ . This gives a  $\mathcal{O}(V(V+E(\log V)))$ .

**Input:** weighted, directed graph  $G = (V,E)$ , with weight function  $w : E \rightarrow \mathbb{R}$ . The weight of path  $p = \langle v_0, v_1, \dots, v_k \rangle$  is the sum of the weights of its constituent edges:

$$w(p) = \sum_{i=1}^k w(v_{i-1}, v_i)$$

The shortest-path weight from  $u$  to  $v$  is

$$\delta(u,v) = \begin{cases} \min\{w(p)\} & \text{if there is path from } u \text{ to } v \\ \infty & \text{otherwise} \end{cases}$$

A shortest path from vertex  $u$  to vertex  $v$  is then defined as any path  $p$  with weight  $w(p) = \delta(u,v)$ .

**All Pairs Shortest Paths:** Compute  $d(u, v)$  the shortest path distance from  $u$  to  $v$  for all pairs of vertices  $u$  and  $v$ .

Assume that the graph is represented by an  $n \times n$  matrix with the weights of the edges.

$$W_{ij} = \begin{cases} 0 & \text{if } i=j \\ w(i,j) & \text{if } i \neq j \text{ \& } (i,j) \in E \\ \infty & \text{if } i \neq j \text{ \& } (i,j) \text{ not } \in E \end{cases}$$

### Floyd-Warshall, Dynamic Programming

- Let  $d^{(k)}_{ij}$  be the weight of a shortest path from vertex  $i$  to vertex  $j$  for which all intermediate vertices are in the set  $\{1, 2, \dots, k\}$ .
- When  $k = 0$ , a path from vertex  $i$  to vertex  $j$  with no intermediate vertex numbered higher than 0 has no intermediate vertices at all, hence  $d^{(0)}_{ij} = w_{ij}$ .

$$d^{(k)}_{ij} = \begin{cases} w_{ij} & \text{if } k=0 \\ \min\{d^{(k-1)}_{ij}, d^{(k-1)}_{ik} + d^{(k-1)}_{kj}\} & \text{if } k \geq 1 \end{cases}$$

## Algorithm

Floyd-Warshall(W)

1  $n \leftarrow \text{rows}[W]$

2  $D^{(0)} \leftarrow W$

3 for  $k \leftarrow 1$  to  $n$

4     do for  $i \leftarrow 1$  to  $n$

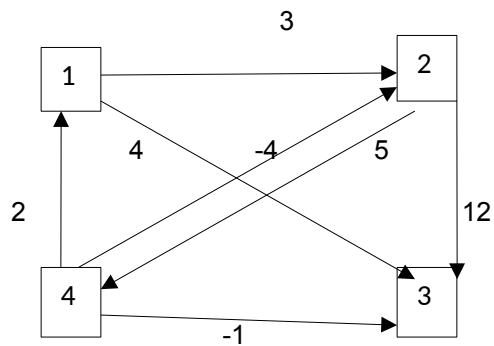
5         do for  $j \leftarrow 1$  to  $n$

6             do  $d^{(k)}_{ij} \leftarrow \min\{d^{(k-1)}_{ij}, d^{(k-1)}_{ik} + d^{(k-1)}_{kj}\}$

7 return  $D^{(n)}$

Running time  $O(V^3)$

Example:



$$D^0 = \begin{pmatrix} 0 & 3 & \infty & \infty \\ \infty & 0 & 12 & 5 \\ 4 & \infty & 0 & -1 \\ 2 & -4 & \infty & 0 \end{pmatrix}$$

$$D^1 \left( \begin{array}{cccc} 0 & 3 & \infty & \infty \\ \infty & 0 & 12 & 5 \\ 4 & 7 & 0 & -1 \\ 2 & -4 & \infty & 0 \end{array} \right) \begin{array}{l} D^1 [2,3] = \min( D^0 [2,3] , D^0 [2,1] + D^0 [1,3] ) = \min(12, \infty + \infty) = 12 \\ D^1 [2,4] = \min( D^0 [2,4] , D^0 [2,1] + D^0 [1,4] ) = \min(5, \infty + \infty) = 5 \\ D^1 [3,2] = \min( D^0 [3,2] , D^0 [3,1] + D^0 [1,2] ) = \min(\infty, 4 + 3) = 7 \\ D^1 [3,4] = \min( D^0 [3,4] , D^0 [3,1] + D^0 [1,4] ) = \min(-1, 4 + \infty) = -1 \\ D^1 [4,2] = \min( D^0 [4,2] , D^0 [4,1] + D^0 [1,2] ) = \min(-4, 2 + 3) = -4 \\ D^1 [4,3] = \min( D^0 [4,3] , D^0 [4,1] + D^0 [1,3] ) = \min(\infty, 2 + \infty) = \infty \end{array}$$

$$D^2 \left( \begin{array}{cccc} 0 & 3 & 15 & 8 \\ \infty & 0 & 12 & 5 \\ 4 & 7 & 0 & -1 \\ 2 & -4 & 8 & 0 \end{array} \right) \begin{array}{l} D^2 [1,3] = \min( D^1 [1,3] , D^1 [1,2] + D^1 [2,3] ) = \min(\infty, 3 + 12) = 15 \\ D^2 [1,4] = \min( D^1 [1,4] , D^1 [1,2] + D^1 [2,4] ) = \min(\infty, 3 + 5) = 8 \\ D^2 [3,1] = \min( D^1 [3,1] , D^1 [3,2] + D^1 [2,1] ) = \min(4, 7 + \infty) = 4 \\ D^2 [3,4] = \min( D^1 [3,4] , D^1 [3,2] + D^1 [2,4] ) = \min(-1, 7 + 5) = -1 \\ D^2 [4,1] = \min( D^1 [4,1] , D^1 [4,2] + D^1 [2,1] ) = \min(2, -4 + \infty) = 2 \\ D^2 [4,3] = \min( D^1 [4,3] , D^1 [4,2] + D^1 [2,3] ) = \min(\infty, -4 + 12) = 8 \end{array}$$

$$D^3 \left( \begin{array}{cccc} 0 & 3 & 15 & 8 \\ 16 & 0 & 12 & 5 \\ 4 & 7 & 0 & -1 \\ 2 & -4 & 8 & 0 \end{array} \right) \begin{array}{l} D^3 [1,2] = \min( D^2 [1,2] , D^2 [1,3] + D^2 [3,2] ) = \min(3, 15 + 7) = 3 \\ D^3 [1,4] = \min( D^2 [1,4] , D^2 [1,3] + D^2 [3,4] ) = \min(8, 15 + (-1)) = 8 \\ D^3 [2,1] = \min( D^2 [2,1] , D^2 [2,3] + D^2 [3,1] ) = \min(\infty, 12 + 4) = 16 \\ D^3 [2,4] = \min( D^2 [2,4] , D^2 [2,3] + D^2 [3,4] ) = \min(5, 12 + (-1)) = 5 \\ D^3 [4,1] = \min( D^2 [4,1] , D^2 [4,3] + D^2 [3,1] ) = \min(2, 8 + 4) = 2 \\ D^3 [4,2] = \min( D^2 [4,2] , D^2 [4,3] + D^2 [3,2] ) = \min(-4, 8 + 7) = -4 \end{array}$$

D<sup>4</sup>

$$\begin{pmatrix} 0 & 3 & 15 & 8 \\ 7 & 0 & 12 & 5 \\ 1 & -5 & 0 & -1 \\ 2 & -4 & 8 & 0 \end{pmatrix}$$

$$D^4 [1,2] = \min( D^3 [1,2] , D^3 [1,4] + D^3 [4,2] ) = \min(3, 8 + (-4)) = 3$$

$$D^4 [1,3] = \min( D^3 [1,3] , D^3 [1,4] + D^3 [4,3] ) = \min(15, 8 + 8) = 15$$

$$D^4 [2,1] = \min( D^3 [2,1] , D^3 [2,4] + D^3 [4,1] ) = \min(16, 5 + 2) = 7$$

$$D^4 [2,3] = \min( D^3 [2,3] , D^3 [2,4] + D^3 [4,3] ) = \min(12, 5 + 8) = 12$$

$$D^4 [3,1] = \min( D^3 [3,1] , D^3 [3,4] + D^3 [4,1] ) = \min(4, -1 + 2) = 1$$

$$D^4 [3,2] = \min( D^3 [3,2] , D^3 [3,4] + D^3 [4,2] ) = \min(7, -1 + (-4)) = -5$$