

Merge sort

1. Split array A[0..n-1] in two about equal halves and make copies of each half in arrays B and C
2. Sort arrays B and C recursively
3. Merge sorted arrays B and C into array A as follows:
 - Repeat the following until no elements remain in one of the arrays:
 - compare the first elements in the remaining unprocessed portions of the arrays
 - copy the smaller of the two into A, while incrementing the index indicating the unprocessed portion of that array
 - Once all elements in one of the arrays are processed, copy the remaining unprocessed elements from the other array into A.

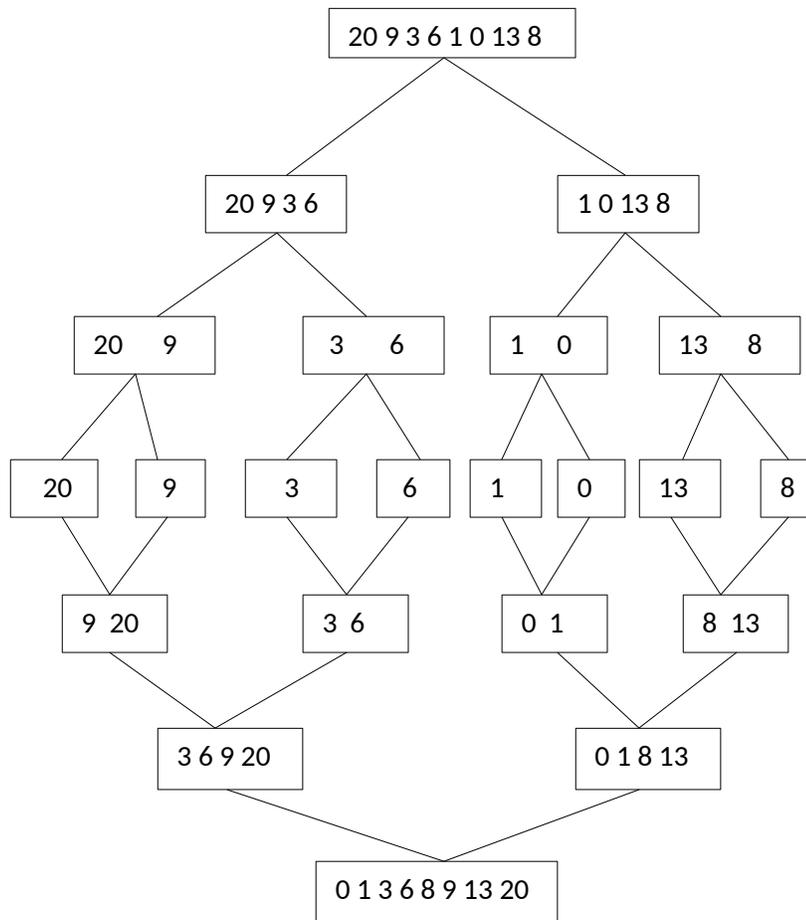
Merge sort algorithm

```
Merge_sort (arr , start_index, last_index )
{
if(start_index== last_index)
    Return(arr[start_index]);
else {
    Mid=floor((start_index+last_index)/2);
    Merge_sort(arr, start_index, mid);
    Merge_sort(arr, mid+1, last_index);
    Merge(arr, start_index, mid, last_index);
}
}
```

Merge algorithm

```
Merge(arr, start_index, mid,last_index)
{ k=mid+1, p=1;
  While(start_index<=mid && k<=last_index)
  {      if(arr[start_index]<arr[k]) {
          Brr[p]=arr[start_index];
          start_index++;
          P++;
        }
    else
    {      Brr[p]=arr[k];
          k++;
          P++;
        }
  }
  for(; start_index<=mid;start_index++)
  {      Brr[p]=arr[start_index]
          P++;
  }
  for(; k<=last_index; k++)
  {      Brr[p]=arr[k];
          P++;
  }
  for(k=1; k<=last_index; k++)
  {
          arr[k]=Brr[k];
  }
}
```

Example of merge sort

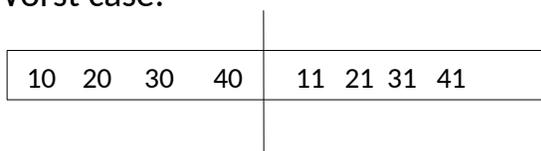


Merge procedure analysis

Input: two sorted sub array

Output: single sorted array

Worst case:



Min(10,11)
 Min(20,11)
 Min(21,20)
 Min(30,21)
 Min(30,31)
 Min(40,31)
 Min(40,41)

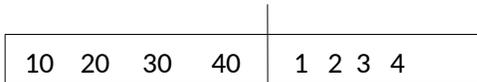


total no. of comparison=(4+4-1)=7 comparison.
worst case time complexity of merge procedure
 $m+n-1 = O(m+n)$ where m & n is the size of sub-array

When both the sub-array size is equal then time complexity

$$n/2+n/2-1=2n=O(n) \quad (\text{neglect constant})$$

Best case:



Min(10,1)
 Min(10,2)
 Min(10,3)
 Min(10,4)



Best case time complexity is no. of comparison
 - if 1 part of array has size m & 2 part of array has size n then time complexity $O(\min(m, n))$
 - If size of sub- arrays is $n/2$ then time complexity $O(n)$

Note: if we don't use second array for merge procedure then time complexity of merge procedure of two sorted sub-array will increase.

Merge sort can be both in-place or outplace but in outplace time complexity is less as compare to in-place because of usage of second array for merge procedure.

Worst case: $O(n)$

Best case: $O(n)$

Average case: $O(n)$

Time complexity :

Merge sort recurrence relation equation

$$T(n) = \begin{cases} O(1) & \text{if } n=1 \\ O(1) + T(n/2) + T(n/2) + O(n) & \text{if } n > 1 \end{cases}$$

Divide cost 2 sub-array conquer cost/ merge cost

$$T(n) = 2T(n/2) + n \dots \text{neglecting constant time}$$

Solved using recurrence relation solving technique. We get time complexity of merge sort **$O(n \log n)$** for outplace sorting.

Space complexity

$$\text{For merge sort : } n + c \log_2 n + n$$

No. of element in array stack size (no. of function calls, where c is no of variable size in each function) array size of b array

space complexity of merge sort = **$O(n)$**

Note: If array size is small then merge sort is not recommended. Merge sort is used for large size array.