

—! Simplex Method :—

Any linear programming problem involving two variables can be easily solved with the help of graphical method as it is easier to deal with two dimensional graph. All the feasible solutions in graphical methods lies within the feasible area on the graph and we used to test the corner points of the feasible solution area for the optimum solution. We used to test all the corner points by putting these values in the objective function.

But if number of variables increase from two, it becomes very difficult to solve the problem by drawing its graph as the problem becomes too complex.

Simplex method is a technique to solve the LPP ~~which~~ that was developed by G.B. Dantzig, an American mathematician. This method is mathematical treatment of the graphical method. Here also various corner points of the feasible area are tested for optimality. But it is not possible to test all the corner points since number of corner points increases manifold as the numbers of equations and variables increases. Therefore, in simplex method the number of corner points to be tested is reduced considerably by using a very effective algorithm which leads us to optimal solution corner point in only a few iterations.

Let us take one example and proceed step by step:

Problem:- 1:

Objective function is to maximise $Z = 12x_1 + 15x_2 + 14x_3$
subjected to conditions

$$-x_1 + x_2 \leq 0 \quad \text{--- ①}$$

$$-x_1 + 2x_3 \leq 0 \quad \text{--- ②}$$

$$x_1 + x_2 + x_3 \leq 100 \quad \text{--- ③}$$

$$x_1, x_2, x_3 \geq 0$$

Solution:-

Step 1:- Right hand side of all the constraints must be either zero or +ve. If they are -ve they must be made +ve by multiplying both side by -1 and sign of inequality would be reversed.

All the inequalities are converted into equalities by adding or subtracting slack or surplus variables. If constraint is \leq , the slack variables are added. If constraint is \geq , then surplus variable is subtracted. In this problem slack variable s_1, s_2, s_3 are added in three equalities:

$$-x_1 + x_2 + s_1 = 0$$

$$-x_1 + 2x_3 + s_2 = 0$$

$$x_1 + x_2 + x_3 + s_3 = 100$$

$$x_1, x_2, s_1, s_2, s_3 \geq 0$$

And objective function becomes

$$\text{maximise } Z = 12x_1 + 15x_2 + 14x_3 + 0 \cdot s_1 + 0 \cdot s_2 + 0 \cdot s_3$$

Step-2 :- We start with a feasible solution and then move towards optimal solution in next iteration.

* Initial feasible solution is preferably chosen to be the origin i.e. x_1, x_2, x_3 assume to be zero.

$$x_1, x_2, x_3 = 0$$

from this we get $s_1 = 0, s_2 = 0$ and $s_3 = 100$. from

Basic variables are the variables which are presently in the solution i.e. s_1, s_2 & s_3 . Non basic variables are the variables which are set to be zero and are not in the current solution e.g. x_1, x_2 & x_3 .

The above information can be expressed in the table

Table 1

| objective function coefficient (C_j) | | 12 | 15 | 14 | 0 | 0 | 0 | |
|--|----------------------|-------|-------|-------|-------|-------|-------|-----|
| e_i | Basic variables (BV) | x_1 | x_2 | x_3 | s_1 | s_2 | s_3 | RHS |
| 0 | s_1 | -1 | 1 | 0 | 1 | 0 | 0 | 0 |
| 0 | s_2 | -1 | 0 | 2 | 0 | 1 | 0 | 0 |
| 0 | s_3 | 1 | 1 | 1 | 0 | 0 | 1 | 100 |

In the above table, 1st row represents coefficients of objective function, 2nd row represent different variables (first regular variables and then slack/surplus). 3rd, 4th and 5th row represents coefficients of variables in all the constraints.

1st column represents coefficients of basic variables in the objective function, 2nd column represent basic variables and last column represent RHS of the constraints in standard form.

In table 1: Current solution is $S_1=0, S_2=0, S_3=100$

Whenever any basic variable assumes zero value, the current solution is said to be degenerate, as in present problem $S_1=0 + S_2=0$ the problem can be further solved by substituting $S_1=t$ and $S_2=t$ where t is very small +ve number.

Step 3: Optimality test:-

Optimality test can be performed to find whether current solution is optimal or not. For this make a new row in the last row in the form of E_j .

Where $E_j = \sum e_i \cdot a_{ij}$

Where a_{ij} - coefficients in the body identity matrix for i^{th} row & j^{th} column.

Now make a new last row $C_j = C_j - E_j$
Table 2.

| | | | | | | | | |
|-------------------------|-------|-------|-------|-------|-------|-------|-------|-----|
| | C_j | 12 | 15 | 14 | 0 | 0 | 0 | |
| e_i | B.V. | x_1 | x_2 | x_3 | S_1 | S_2 | S_3 | RHS |
| 0 | S_1 | -1 | 1 | 0 | 1 | 0 | 0 | t |
| 0 | S_2 | -1 | 0 | 2 | 0 | 1 | 0 | t |
| 0 | S_3 | 1 | 1 | 1 | 0 | 0 | 1 | 100 |
| $E_j = \sum e_i a_{ij}$ | | 0 | 0 | 0 | 0 | 0 | 0 | |
| $C_j = C_j - E_j$ | | 12 | 15 | 14 | 0 | 0 | 0 | |

If any of the value of $(C_j - E_j)$ is +ve, then it means that most +ve value will give the key column. In the present case x_2 is potential variable (most +ve value).

If all the values in $(C_j - E_j)$ is -ve then it means optimum solution has been reached.

⇒ Step 4 :- Iterate towards optimum solution :-

Maximum value of $C_j - E_j$ gives the key column.

Now find out key row. For this first introduce ratio column at the last. To obtain the ratio column, divide the coefficients in RHS column by the corresponding coefficients in the key column.

Now look for the least +ve value in the ratio column and that would give the key row.

| | C_j | 12 | 15 | 14 | 0 | 0 | 0 | | |
|-------|-------------|-------|-------|-------|-------|-------|-------|-----|-------------|
| E_i | B.V. | x_1 | x_2 | x_3 | S_1 | S_2 | S_3 | RHS | Ratio |
| 0 | S_1 | -1 | [1] | 0 | 1 | 0 | 0 | t | t → Key row |
| 0 | S_2 | -1 | 0 | 2 | 0 | 1 | 0 | t | 1/2 (∞) |
| 0 | S_3 | 1 | 1 | 1 | 0 | 0 | 1 | 100 | 100 |
| | E_j | 0 | 0 | 0 | 0 | 0 | 0 | | |
| | $C_j - E_j$ | 12 | 15 | 14 | 0 | 0 | 0 | | |

↑
Key column

In the present problem we have three values in ratio column i.e. t, 1/2, & 100 out of these t is the least +ve. Therefore row corresponding to t in ratio column would be the key row.

Key column decides entering variable and key row decides leaving variable. Therefore in the present problem x_2 is entering variable i.e. it would become the basic and would enter in the solution. And S_1 would be the leaving variable.

* Element at the intersection of key row & key column is the key element.

Now all the element in the key column except the key element is to be made zero and key element is to be made unity.

Table 3

| | | | | | | | | | |
|-------|-------------|-------|-------|-------|-------|-------|-------|---------|---|
| | C_j | 12 | 15 | 14 | 0 | 0 | 0 | | |
| e_i | BV | x_1 | x_2 | x_3 | s_1 | s_2 | s_3 | RHS | |
| 15 | x_2 | -1 | 1 | 0 | 1 | 0 | 0 | t | |
| 0 | s_2 | -1 | 0 | 2 | 0 | 1 | 0 | t | |
| 0 | s_3 | 2 | 0 | 1 | -1 | 0 | 0 | $100-t$ | $\rightarrow R_3 \rightarrow R_3 - R_1$ |
| | E_j | -15 | 15 | 0 | 15 | 0 | 0 | | |
| | $C_j - E_j$ | 27 | 0 | 14 | -15 | 0 | 0 | | |

In the above table s_1 is replaced by x_2 in the basic variables and corresponding e_i column has also been changed.

* Repeat the solution iteration until the optimal solution is obtained i.e. all elements in the $C_j - E_j$ row becomes zero or -ve.

| | | | | | | | | | |
|-------|-------------|------------|-------|-------|-------|-------|-------|---------|--|
| | C_j | 12 | 15 | 14 | 0 | 0 | 0 | | |
| e_i | BV | x_1 | x_2 | x_3 | s_1 | s_2 | s_3 | RHS | Ratio |
| 15 | x_2 | -1 | 1 | 0 | 1 | 0 | 0 | t | -t |
| 0 | s_2 | -1 | 0 | 2 | 0 | 1 | 0 | t | -t |
| 0 | s_3 | $[2]^{KE}$ | 0 | 1 | -1 | 0 | 1 | $100-t$ | $50 - \frac{t}{2} \leftarrow \text{key row}$ |
| | E_j | -15 | 15 | 0 | 15 | 0 | 0 | | |
| | $C_j - E_j$ | 27 | 0 | 14 | -15 | 0 | 0 | | |
| | | \uparrow | | | | | | | |
| | | key column | | | | | | | |

Now x_1 is entering value and S_3 is leaving value. Make key element unity and all the key column element other than key element zero.

| | | | | | | | | | | |
|-------|-------------|-------|---------------|----------------------|----------------|-------|-----------------|--------------------|-------------------------------|------------------------------|
| | C_j | | 12 | 15 | 14 | 0 | 0 | 0 | | |
| e_i | BV | x_1 | x_2 | x_3 | s_1 | s_2 | s_3 | RHS | | Ratio |
| 15 | x_2 | 0 | 1 | $\frac{1}{2}$ | $\frac{1}{2}$ | 0 | $\frac{1}{2}$ | $50 + \frac{t}{2}$ | $(R_1 \rightarrow R_1 + R_3)$ | $100 + t$ |
| 0 | s_2 | 0 | 0 | $[\frac{5}{2}]_{KF}$ | $-\frac{1}{2}$ | 1 | $\frac{1}{2}$ | $50 + \frac{t}{2}$ | $(R_2 \rightarrow R_2 + R_3)$ | $20 + \frac{t}{5}$ ← key row |
| 12 | x_1 | 1 | 0 | $\frac{1}{2}$ | $-\frac{1}{2}$ | 0 | $\frac{1}{2}$ | $50 - \frac{t}{2}$ | | $100 - t$ |
| | E_j | 12 | 15 | $\frac{27}{2}$ | $\frac{3}{2}$ | 0 | $\frac{27}{2}$ | | | |
| | $C_j - E_j$ | 0 | 0 | $\frac{1}{2}$ | $-\frac{3}{2}$ | 0 | $-\frac{27}{2}$ | | | |

↑
Key Column

| | | | | | | | | | |
|-------|-------------|-------|-------|-------|----------------|----------------|-----------------|---------------------|---|
| | C_j | | 12 | 15 | 14 | 0 | 0 | 0 | |
| e_i | BV | x_1 | x_2 | x_3 | s_1 | s_2 | s_3 | RHS | |
| 15 | x_2 | 0 | 1 | 0 | $\frac{3}{5}$ | $-\frac{1}{5}$ | $\frac{2}{5}$ | $40 + \frac{2t}{5}$ | $(R_1 \rightarrow R_1 - \frac{R_2}{2})$ |
| 14 | x_3 | 0 | 0 | 1 | $-\frac{1}{5}$ | $\frac{2}{5}$ | $\frac{1}{5}$ | $20 + \frac{t}{5}$ | |
| 12 | x_1 | 1 | 0 | 0 | $-\frac{2}{5}$ | $-\frac{1}{5}$ | $\frac{2}{5}$ | $40 - \frac{3t}{5}$ | $(R_3 \rightarrow R_3 - \frac{R_2}{2})$ |
| | E_j | 12 | 15 | 14 | $+\frac{7}{5}$ | $\frac{1}{5}$ | $\frac{68}{5}$ | | |
| | $C_j - E_j$ | 0 | 0 | 0 | $-\frac{7}{5}$ | $-\frac{1}{5}$ | $-\frac{68}{5}$ | | |

It can be seen that since all the values in $C_j - E_j$ row are either -ve or zero, hence optimal solution has been reached.

Final solution $x_1 = 40 - \frac{3t}{5}$; t is very small quantity and hence can be neglected

therefore

$x_1 = 40$
 $x_2 = 40$
 $x_3 = 20$